# Discovering Hidden Structures in Complex Networks

## Roman Vershynin







March 2015

# Network Science is highly interdisciplinary.



+ finance + technology  $+ \dots$ 

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Some structures are apparent, local.





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Protein interaction network [A.-L. Barabási & Z. Oltvai, Nature Reviews Genetics 5, 101–113, Feb. 2004]

Some structures are apparent, local.



#### The Internet

(C. Hurter et al., Eurographics Conference on Visualization 2012) (C. Hurter et al., Eurographics Conference on Visualization 2012)

Other structures are latent, global...



#### Collaboration network of economists

(AER, JPE, Econometrica, RES, QJE. www.cloudycnénfinet) ( 豆 ト ( 豆 ト ) 豆 のの(

... just like in nature:



#### global structure

local chaos

## **Basic** Questions

- How can we find latent structures in real networks?
- How can we explain and model these structures?

## Mathematical perspective

Model large networks as random graphs.

A leap of faith.

(Edges drawn at random.)

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# Random graphs: Erdös-Rényi model G(n, p)

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G(n, p) with n = 1000, p = 0.00095



# Inhomogeneous Erdös-Rényi model $G(n, (p_{ij}))$

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Example. Stochastic block model with two communities G(n, p, q): Edges within each community: probability p; across communities: probability q < p.





# Inhomogeneous Erdös-Rényi model $G(n, (p_{ij}))$

Multiple communities are possible to model, too:

Stochastic block model

Real data (aggression network of students)





(UC Davis Center for Visualization)

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# Network Model Recovery

Model Recovery Problem. Observe one instance of a network from  $G(n, (p_{ij}))$ . Recover the model, i.e. the connection probabilities  $p_{ij}$ .

Application to real graphs:



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#### Application to real graphs:



 $p_{ij} =$  "latent bonds" between vertices.

Link prediction.

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# Network Model Recovery Problem

A particular case, for stochastic block models:

Community Detection Problem. Observe a network drawn from the stochastic block model G(n, p, q). Recover the two communities.



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### From graphs to matrices

#### Adjacency matrix A:



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For inhomogeneous Erdös-Rényi model:

 $A = (\text{Bernoulli}(p_{ij}))$   $\mathbb{E} A = (p_{ij})$ 

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Model Recovery Problem. Observe A; recover  $\mathbb{E} A$ .

## Relation to matrix completion

Evident but not thoroughly explored.

Matrix completion: recover a low-rank matrix from a few randomly chosen entries.



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Network model recovery: recover a (low-rank?) matrix  $\mathbb{E} A = (p_{ij})$  from random measurements  $A = (\text{Bernoulli}(p_{ij}))$ .

<b>F</b> 0	1	0	1	0	0	0	07		<b>F</b> 1	.7	.6	.7	.1	.4	.3	.2
1	0	0	0	0	0	1	1		.7	1	.6	.5	.2	.1	.2	.1
0	0	0	0	1	1	0	0		.6	.6	1	.9	.4	.2	.3	.3
1	0	0	0	0	0	0	1	?	.7	.5	.9	1	.2	.1	.3	.2
0	0	1	0	0	0	1	0	$\rightarrow$	.1	.2	.4	.2	1	.8	.6	.5
0	0	1	0	0	1	0	0		.4	.1	.2	.1	.8	1	.7	.6
0	1	0	0	0	0	0	1		.3	.2	.3	.3	.6	.7	1	.9
LO	1	0	1	0	0	1	0		L.2	.1	.3	.2	.5	.6	.9	1

Most relevant comparison is to single-bit matrix completion [Davenport et al '12].

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# Existing approaches

Mostly apply to stochastic block models.

Insights from Combinatorics, Computer Science, Statistics, Physics:

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- combinatorial techniques (min-cut, hierarchical clustering)
- spectral methods this talk
- statistical inference (likelihood maximization)
- variational methods
- Markov chain Monte Carlo
- belief propagation
- convex optimization
- semidefinite programming this talk
- . . .

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Is this true? In other words:

Question. Do random graphs concentrate near their "expected" graphs?

Consider an inhomogeneous Erdös-Rényi random graph  $G(n, (p_{ij}))$  with expected degrees  $np_{ij} \sim d$ .



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Theorem. A random graph with expected degrees  $d \ge \log n$  concentrates:

 $\|A - \mathbb{E} A\| \lesssim \sqrt{d} \quad w.h.p. \ while \quad \|\mathbb{E} A\| \sim d.$ 



Consider an inhomogeneous Erdös-Rényi random graph  $G(n, (p_{ij}))$  with expected degrees  $np_{ij} \sim d$ .



Theorem. A random graph with expected degrees  $d \gtrsim \log n$  concentrates:

 $\|A-\mathbb{E}\,A\|\lesssim \sqrt{d}\quad w.h.p. \ while \quad \|\,\mathbb{E}\,A\|\sim d.$ 

#### Proofs:

• [Kahn-Szemeredi 89]  $\rightarrow$  [Feige-Ofek 05, Lei-Rinaldo 13, Chin-Rao-Vu 15]: Simple concentration of  $x^{\mathsf{T}}(A - \mathbb{E}A)y$  for fixed x, y; then complicated union bound over x, y (tailored the coefficient profiles of x, y).

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- Other approaches: [Hajek-Wu-Xu 14; Bandeira-van Handel 14; Le-Vershynin 15].
- Weaker results: [Furedi-Komlos 80] with  $d \gtrsim \log^4 n$ ; [Oliveira 10] with  $||A - \mathbb{E}A|| \lesssim \sqrt{d \log n}$  by matrix Bernstein inequality.

### Sparse random graphs do not concentrate

Observation. A random graph G(n, p) with expected degrees  $d = np \ll \log n$  does not concentrate:

 $||A - \mathbb{E}A|| \gg ||\mathbb{E}A||.$ 

See [Krivelevich-Sudakov 03].



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What is wrong with sparse graphs?

The degrees are wild, do not concentrate near d anymore. **High-degree vertices** blow up ||A||: some columns of A are too large.

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## Sparse random graphs do not concentrate

High-degree vertices dominate the picture. Spectral methods reveal only those vertices. *Local information*, no latent structure [Mihail-Papadimitriou 02].



Preprocess the network.

Regularize the high-degree vertices: reweight (or remove) enough edges from them.

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• Yes, if we remove all high-degree vertices and all their edges [Feige-Ofek 05]. But these vertices hold the network together (hubs)! Their removal can cause network to fall apart.

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- Yes, if we remove all high-degree vertices and all their edges [Feige-Ofek 05]. But these vertices hold the network together (hubs)! Their removal can cause network to fall apart.
- Yes, in full generality. Any type of regularization helps, as long as it brings down the degrees to  $\sim d$ . [Le-Levina-V, Le-V 05].

Inhomogeneous E-R random graph with  $d = \max np_{ij}$ .

Regularize vertices with degrees > 2d: make all degrees  $\leq 2d$  by reducing the weights of edges arbitrarily.

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#### Proof:

- simple concentration of A in cut norm;
- **2** upgrade to operator norm on a subgraph by *Grothendieck-Pietsch factorization*;

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iteration to extend the control over all graph.

By-product: a new graph decomposition.

## Regularization and concentration: applications

Eigenvectors reveal the latent structure?



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Eigenvectors reveal the latent structure?

Concentration (possibly after regularization)  $\Rightarrow$ 

 $A\approx \mathbb{E}\,A.$ 

Davis-Kahan theorem  $\Rightarrow$  eigenvectors satisfy

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Eigenvectors  $v_i(\mathbb{E}A)$  carry information about **network structure**.

Example. Community detection in stochastic block model G(n, p, q).





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Example. Community detection in stochastic block model G(n, p, q).



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Example. Community detection in stochastic block model G(n, p, q).



 $v_2(\mathbb{E} A)$  encodes community structure  $\Rightarrow v_2(A)$  encodes the structure, too.

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Example. Community detection in stochastic block model G(n, p, q).



 $v_2(\mathbb{E} A)$  encodes community structure  $\Rightarrow v_2(A)$  encodes the structure, too.

Spectral Clustering Algorithm: given a graph with adjacency matrix A,

- Compute the second leading *eigenvector* of *A*;
- Recover communities based on the signs of its coefficients.

## Using eigenvectors: theory.

Corollary (Community Detection). Consider the stochastic block model G(n, p, q) with p = a/n and q = b/n. Suppose

 $(a-b)^2 \ge C_{\varepsilon}(a+b).$ 

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Proof: straightforward consequence of concentration [Le-Levina-V; Le-V 05].

Detection threshold. The condition on is *optimal* up to  $C_{\varepsilon}$ , which must  $\to \infty$ . No algorithm can succeed if

 $(a-b)^2 \le 2(a+b).$ 

There are algorithms that do better than random guess if

 $(a-b)^2 > 2(a+b).$ 

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See [Mossel-Neeman-Sly 13-14; Massoulié 13; Bordenave-Lelarge-Massoulié 15].

## Performance of regularized spectral clustering



n = 400 vertices, expected degree 5. Connection probabilities p = 5/n and b = 0.5/n.

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Further application of

 $\operatorname{eigenstructure}(A) \approx \operatorname{eigenstructure}(\mathbb{E} A).$ 



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eigenstructure(A)  $\approx$  eigenstructure( $\mathbb{E} A$ ).

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Assume  $\mathbb{E} A$  has low rank, exactly or approximately. Then PCA on A should **reveal the latent structure** of the network.

How? project the columns of A onto the space of the 3 leading eigenvectors.

Power grid of U.S.A.



Without regularization:



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Not very useful...

With regularization:



SOR

With regularization:



SAG

With regularization:



na a

With regularization:



SOC

With regularization:



DQ CP

With regularization:



Sac

With regularization:



SOR

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SAR

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Diffusion approach: heat the graph.

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Diffusion approach: heat the graph.

The heat gets trapped in a community  $\Rightarrow$  can recover it.



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In  $\mathbb{R}^2$ , the **heat diffusion** is described by the Laplacian  $\Delta = \frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2}$ .



From Gabriel Peyré's manifold methods class (left); Morpheo research team (right)



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On a graph, the discrete Laplacian is the  $n \times n$  matrix

 $\Delta := I - D^{-1/2} A D^{-1/2}$ 

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where D is the diagonal matrix with the *degrees* on the diagonal.

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Adjacency and Laplacian are two most fundamental matrices associated to graphs.

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For sparse graphs  $(d \ll \log n)$ , fails to concentrate.

For **dense graphs** (expected degrees  $d \gtrsim \log n$ ), Laplacian concentrates.

For sparse graphs  $(d \ll \log n)$ , fails to concentrate.

What's wrong? Low-degree vertices: isolated vertices, trees. (They get overheated.)



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Would regularization help?

#### Would regularization help?

**Connect** low-degree vertices to the rest of the graph by *light weighted edges*; bring up all degrees to  $\sim d$ .



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Proposed by network scientists [Chaudhuri+ 12, Amini+ 13].

Theorem. The Laplacian  $\Delta'$  of the regularized graph concentrates:

 $\|\Delta' - \mathbb{E} \Delta'\| \lesssim \frac{1}{\sqrt{d}} \quad while \quad \|\Delta'\| \sim 1.$ 

[Le-Levina-V, Le-V 05].



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Application to community detection: use the  $2^{nd}$  eigenvector of the Laplacian. Theoretical performance: same as for adjacency; empirically even better.

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Physical interpretation: Make the graph vibrate; the wave with lowest frequency recovers the communities.



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# Performance of regularized spectral clustering

Artificial data: sparse stochastic block model

Without regularization

With regularization



This tree gets overheated



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# Performance of regularized spectral clustering

Real data: political blogs after 2004 U.S. presidential election [Adamic-Glance 04].

With regularization

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1,222 vertices (liberal/conservative); edges = hyperlinks; average degree = 27.

Goal: fit the desired type of structure to a given network.

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Strongest community structure: union of cliques.

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Optimization:  $\max \langle A, Z \rangle$  where A = adjacency matrix of the network, Z = adjacency matrix of a union of cliques with k edges.



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Optimization:  $\max \langle A, Z \rangle$  where A = adjacency matrix of the network, Z = adjacency matrix of a union of cliques with k edges.



Integer optimization problem. NP-hard.

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### Semidefinite relaxation

**Fact.** A matrix  $Z \in \{0, 1\}^{n \times n}$  is block diagonal  $\Leftrightarrow Z$  is positive semidefinite.

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A semidefinite relaxation:

SDP.  $\max \langle A, Z \rangle$ :  $Z \in [0, 1]^{n \times n}$  is positive semidefinite,  $\sum Z_{ij} = k$ .

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General stochastic block model:  $\forall$  many communities,  $\forall$  connection probabilities  $p_{ij}$ , within communities > p; across communities < q. (Not necessarily low rank!)



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Theorem (Community Detection by SDP). Consider a general stochastic block model with p = a/n and q = b/n. Suppose

 $(a-b)^2 \ge C_{\varepsilon}(a+b).$ 

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Then the SDP (with k=number of edges) recovers communities up to  $\varepsilon n$  misclassified vertices, and with high probability.

[Guedon-V. 14].

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Exact recovery for dense networks  $(a, b \ge \log n)$ ; thresholds known [Abbeet al  $\ge 14$ ]  $\circ \circ \circ$ 

Example. Dolphins in Doubtful Sound, New Zealand [Lusseau et al. 03].



True communities



Communities found by SDP



Take a closer look at

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**Output:** *k* strongest "latent bonds" between vertices.

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## Semidefinite relaxation in action

Take a closer look at

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**Output:** *k* strongest "latent bonds" between vertices.



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**Next slide:** increase k gradually  $\Rightarrow$  dynamic picture.



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## Performance of semidefinite relaxation

SDP enhances the latent structure of the network:



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SDP densifies communities, sparsifies cuts across communities.

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SDP densifies communities, sparsifies cuts across communities.

SDP did not know the number of communities in advance. It decided that 2 communities should fit best.

## Compressed sensing

Signal: vector, matrix Structure: sparsity, low rank Measurements: random linear, few Outliers: permitted in robust PCA Exact recovery; exact thresholds Recent blowup (2004+)

## Structure recovery in networks

Signal: network model  $(p_{ij})$ Structure: low rank, ??? (open) Measurements: 0/1 random, few Outliers: permitted (high/low degree vertices) Exact recovery; exact thresholds Recent blowup (2012+)

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